# On the difficulty of inverting automorphisms of free groups 

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## Outline

(1) Motivation
(2) Free groups
(3) Lower bounds: a good enough example

4 Upper bounds: outer space
(5) The special case of rank 2

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4 Upper bounds: outer space
(5) The special case of rank 2

# (Joint work with P. Silva and M. Ladra.) 

Find a group G where • is "easy" but ( $)^{-1}$ is "difficult"
Natural candidate: Aut $\left(F_{n}\right)$, where $F_{r}=\left\langle a_{1}\right.$
$F_{3}=\langle a, b, c \mid\rangle$.


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\begin{aligned}
\phi \psi: F_{3} & \rightarrow F_{3} \\
a & \mapsto b c^{-1} a^{-1} b c \\
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c & \mapsto a^{-1} b c^{-1} .
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F_{5}=\langle a, b, c, d & , & & \\
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## - We have formalized the situation.

- We have seen that inverting in $\operatorname{Aut}\left(F_{r}\right)$ is not that bad.
- We now want to look for worse groups $G$.


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## Main definition

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Let $A=\left\{a_{1}, \ldots, a_{r}\right\}$ be a finite alphabet, and $G=\langle A \mid R\rangle$ be a finite presentation for a group $G$. We have the word metric:

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\text { for } g \in G, \quad|g|=\min \left\{n \mid g=a_{i_{1}}^{\epsilon_{1}} \cdots a_{i_{n}}^{\epsilon_{n}}\right\} .
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## Definition

For $\theta \in \operatorname{Aut}(G)$, note $\theta$ is determined by $a_{1} \theta, \ldots, a_{r} \theta$ and define

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\|\theta\|_{\infty}=\max \left\{\left|a_{1} \theta\right|, \ldots,\left|a_{r} \theta\right|\right\} .
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For every $\theta \in A$

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\alpha_{A}(n)=\max \left\{\left\|\theta^{-1}\right\|_{1} \mid \theta \in \operatorname{Aut}(G),\|\theta\|_{1} \leqslant n\right\} .
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Clearly, $\alpha_{A}(n) \leqslant \alpha_{A}(n+1)$.

The bigger is $\alpha_{A}$, the more "difficult" will be to invert automorphisms of $G$ (with respect to the given set of generators $A$ ).

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Determine the asymptotic growth of the function $\alpha_{A}$.

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## Free group case

For the rest of the talk, $G=F_{r}=\left\langle a_{1}, \ldots, a_{r} \mid\right\rangle$.

## Definition

Every $w \in F_{r}$ has its length, $|w|$, and its cyclic length, $|w|$
$\left|a_{1} a_{1}^{-1} a_{2}\right|=\left|a_{2}\right|=\left|a_{2}\right|=1$,
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i) $\left|w^{n}\right| \leqslant|n||w|$ and $\cdot\left|w^{n}\right| \cdot=|n| \cdot|w| \cdot$
ii) $|v w| \leqslant|v|+|w|$, but $\cdot|v w| \cdot \leqslant|v| \cdot+|w| \cdot$ is not true in general.

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## but not equal in general.

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Consider $\theta: F_{4} \rightarrow F_{4}, a \mapsto a, b \mapsto a^{-1} b a, c \mapsto a^{-1} c a, d \mapsto d$. We have $\left|\cdot \theta\left\|_{1}=4,\right\| \theta\right| \|_{1}=6$ and $\|\theta\|_{1}=8$.

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Are these functions equal up to multiplicative constants ?
$\alpha_{r}$ and $\gamma_{r}$ are not;
$\beta_{r}$ is not clear.

## Main results

## Theorem

For rank $r=2$ we have
(i) for $n \geqslant 4, \quad \alpha_{2}(n) \leqslant \frac{(n-1)^{2}}{2}$,
(ii) for $n \geqslant n_{0}, \quad \alpha_{2}(n) \geqslant \frac{n^{2}}{16}$,
(iii) for $n \geqslant 1, \beta_{2}(n)=n$,
(iv) for $n \geqslant 1 . \gamma_{2}(n)=n$.

## Theorem

For $r \geqslant 3$ there exist $K=K(r)$ and $M=M(r)$ such that, for $n \geqslant 1$,
(i) $\alpha_{r}(n) \geqslant K n^{r}$,
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## A lower bound for $\gamma_{r}$

## Theorem

For $r \geqslant 2$, and $n \geqslant n_{0}$, we have $\gamma_{r}(n) \geqslant \frac{1}{2 r^{r-1}} n^{r-1}$.
Proof: For $r \geqslant 2$ and $n \geqslant 1$, consider

| $\psi_{r, n}: F_{r}$ | $\rightarrow F_{r}$ | $\psi_{r, n}^{-1}: F_{r}$ | $\rightarrow$ | $F_{r}$ |
| ---: | :--- | ---: | :--- | :--- |
| $a_{1}$ | $\mapsto a_{1}$ | $a_{1}$ | $\mapsto$ | $a_{1}$ |
| $a_{2}$ | $\mapsto a_{1}^{n} a_{2}$ | $a_{2}$ | $\mapsto$ | $a_{1}^{-n} a_{2}$ |
| $a_{3}$ | $\mapsto a_{2}^{n} a_{3}$ |  | $\vdots$ |  |

$$
a_{r} \mapsto a_{r-1}^{n} a_{r}
$$



A straightforward calculation shows that
$\left\|\psi_{r, n}\right\|_{1}=\left\|\psi_{r, n}\right\|_{1}=(r-1) n+r$, and
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For $r \geqslant 2$, and $n \geqslant n_{0}$, we have $\gamma_{r}(n) \geqslant \frac{1}{2 r^{r-1}} n^{r-1}$.
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\begin{array}{rlrll}
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a_{3} & \mapsto & a_{2}^{n} a_{3} & & a_{2} \\
& \vdots & & a_{1}^{-n} a_{2} \\
a_{r} & \mapsto & a_{r-1}^{n} a_{r} & & a_{i} \\
& & \mapsto & \\
& & & (2 \leqslant i \leqslant r) &
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Hence, for $n \geqslant r$,

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## Now, for $n$ big enough, take the closest multiple of $r$ below,



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## Outline

## (1) <br> Motivation

(2) Free groups

3 Lower bounds: a good enough example

4 Upper bounds: outer space
(5) The special case of rank 2

## Outer space

To prove the upper bound
(ii) $\beta_{r}(n) \leqslant K n^{M}$,
we'll need to use the recently discovered metric in the outer space $\mathcal{X}_{r}$.

## Definition

- By graf $\Gamma$ we mean a finite, connected graph of rank $r$, with no vertices of degree 1 or 2.
- A metric on $\Gamma$ is a map $\ell: E \Gamma \rightarrow[0,1]$ such that $\sum_{e \in E \Gamma} \ell(e)=1$, and $\{e \in E \Gamma \mid \ell(e)=0\}$ is a forest.
- For a graph $\Gamma, \Sigma_{\Gamma}=\{$ metrics on $\Gamma\}=$ a simplex with missing faces.
- If $\Gamma^{\prime}=\Gamma /$ forest, then we identify points in $\Sigma_{\Gamma^{\prime}}$ with the corresponding points in $\Sigma_{\Gamma}$ by assigning length 0 to the collapsed ed'ges.
- A marking on $\Gamma$ is a homotopy equivalence $f: R_{r} \rightarrow \Gamma$.


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The outer space $\mathcal{X}_{r}$ is

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(where $\sim$ is an equivalence relation).

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There is a natural action of $\operatorname{Aut}\left(F_{r}\right)$ on $\mathcal{X}_{r}$, given by
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## Metric on $\mathcal{X}_{r}$

## Definition

Let $x, x^{\prime} \in \mathcal{X}_{r}, x=(\Gamma, f, \ell), x^{\prime}=\left(\Gamma^{\prime}, f^{\prime}, \ell^{\prime}\right)$. A difference of markings is a map $\alpha: \Gamma \rightarrow \Gamma^{\prime}$, which is linear over edges and $f \alpha \simeq f^{\prime}$.
For such an $\alpha$, define $\sigma(\alpha)$ to be its maximum slope over edges.

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$\mathcal{X}_{r}$ admits the following "metric":
$d\left(x, x^{\prime}\right)=\min \{\log (\sigma(\alpha)) \mid \alpha$ diff. markings $\}$
This minimum is achieved by Arzela-Ascoli's theorem.
This is Bestvina-AlgomKfir version of Martino-Francaviglia's original metric.

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## Proposition

(i) $d(x, y) \geqslant 0$, and $=0 \Leftrightarrow x=y$.

$$
\text { (ii) } d(x, z) \leqslant d(x, y)+d(y, z) \text {. }
$$

(iii) $\operatorname{Out}\left(F_{r}\right)$ acts by isometries, i.e. $d(\phi \cdot x, \phi \cdot y)=d(x, y)$.
(iv) But... $d(x, y) \neq d(y, x)$ in general.

## Definition

For $\epsilon>0$, the $\epsilon$-thick part of $\mathcal{X}_{r}$ is

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\mathcal{X}_{r}(\epsilon)=\left\{(\Gamma, f, \ell) \in \mathcal{X}_{r} \mid \ell(p) \geqslant \epsilon \forall \text { closed path } p \neq 1\right\}
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## Bestvina-AlgomKfir theorem

## Theorem (Bestvina-AlgomKfir)

For any $\epsilon>0$ there is constant $M=M(r, \epsilon)$ such that for all $x, y \in \mathcal{X}_{r}(\epsilon)$,

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## Corollary

For $r \geqslant 2$, there exists $M=M(r)$ such that

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Corollary
For $r \geqslant 2$, there exists $M=M(r)$ such that

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\beta_{r}(n) \leqslant r n^{M} .
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## Proof

Remind $\beta_{r}(n)=\max \left\{\left|\left\|\theta^{-1}\left|\|_{1}\right| \theta \in\right.\right.\right.$ Aut $\left.\left.\left.F_{r},\right\|\right|\|\theta\|_{1} \leqslant n\right\}$.
Proof. Given $\phi \in \operatorname{Aut}\left(F_{r}\right)$, consider $x=\left(R_{r}, i d, \ell_{0}\right) \in \mathcal{X}_{r}$, and $\phi \cdot x=\left(R_{r}, \phi, \ell_{0}\right) \in \mathcal{X}_{r}$, where $\ell_{0}$ is the uniform metric.
$d(x, \phi \cdot x)=\min \{\log (\sigma(\alpha)) \mid \alpha$ diff. markings $\}$


Now, using Bestvina-AlgomKfir theorem,
$\log \left(\left|\left\|\phi^{-1} \mid\right\|_{1}\right) \sim d\left(x, \phi^{-1} \cdot x\right)=d(\phi \cdot x, x) \leq M d(x, \phi \cdot x) \sim M \log \left(\| \| \phi \|_{1}\right)\right.$
Hence, for every $\phi \in \operatorname{Aut}\left(F_{r}\right),\| \| \phi^{-1}\| \|_{1} \leqslant r\| \| \phi \|_{1}^{M} . \square$

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Now, using Bestvina-AlgomKfir theorem,


## Proof

Remind $\beta_{r}(n)=\max \left\{\left|\left\|\theta^{-1}\left|\|_{1}\right| \theta \in\right.\right.\right.$ Aut $\left.\left.F_{r},\right\|\|\theta\| \|_{1} \leqslant n\right\}$.
Proof. Given $\phi \in \operatorname{Aut}\left(F_{r}\right)$, consider $x=\left(R_{r}, i d, \ell_{0}\right) \in \mathcal{X}_{r}$, and $\phi \cdot x=\left(R_{r}, \phi, \ell_{0}\right) \in \mathcal{X}_{r}$, where $\ell_{0}$ is the uniform metric.

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\begin{aligned}
d(x, \phi \cdot x) & =\min \{\log (\sigma(\alpha)) \mid \alpha \text { diff. markings }\} \\
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$\log \left(\left\|\phi^{-1} \mid\right\|_{1}\right) \sim d\left(x, \phi^{-1} \cdot x\right)=d(\phi \cdot x, x) \leqslant M d(x, \phi \cdot x) \sim M \log \left(\left|\|\phi \mid\|_{1}\right)\right.$.
Hence, for every $\phi \in \operatorname{Aut}\left(F_{r}\right),\left\|\left|\phi^{-1}\right|\right\|_{1} \leqslant r\|\mid \phi\|_{1}^{M}$. $\square$

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## Outline

## (1) Motivation

(2) Free groups

3 Lower bounds: a good enough example

4 Upper bounds: outer space
(5) The special case of rank 2

## The rank 2 case

These functions for Aut $\left(F_{2}\right)$ are much easier to understand due to the following technical lemmas.

Lemma
Let $\varphi \in \operatorname{Aut}\left(F_{2}\right)$ be positive. Then $\varphi^{-1}$ is cyclically reduced and

## Lemma

For everv $\theta \in \operatorname{Aut}\left(F_{2}\right)$, there exist two letter permuting autos $\psi_{1}, \psi_{2} \in \operatorname{Aut}\left(F_{2}\right)$, a positive one $\varphi \in \operatorname{Aut}^{+}\left(F_{2}\right)$, and an element $g \in$

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## The rank 2 case: $\gamma_{2}$

Theorem
For every $\theta \in \operatorname{Aut}\left(F_{2}\right),\left\|\cdot \theta^{-1}\right\|_{1}=H \theta \|_{1}$. Hence, $\gamma_{2}(n)=n$.

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For $n \geqslant 4$ we have $\alpha_{2}(n) \leqslant \frac{(n-1)^{2}}{2}$.
Proof. Let $\theta \in \operatorname{Aut}\left(F_{2}\right)$, decomposed as above, $\theta=\psi_{1} \varphi \psi_{2} \lambda_{g}$. Then, $\theta^{-1}=\lambda_{g^{-1}} \psi_{2}^{-1} \varphi^{-1} \psi_{1}^{-1}$ and

$$
\left\|\theta^{-1}\right\|_{1} \leqslant 4|g| \cdot\left\|\psi_{2}^{-1} \varphi^{-1} \psi_{1}^{-1}\right\|_{\infty}=4|g| \cdot\left\|\varphi^{-1}\right\| \infty
$$

$$
4|g|\left(\left\|\varphi^{-1}\right\|_{1}-1\right)=4|g|\left(\|\varphi\|_{1}-1\right) .
$$

Now from $\|\varphi\|_{1}+2|g| \leqslant\|\theta\|_{1}=n$, we deduce $|g| \leqslant \frac{n-\|\varphi\|_{1}}{2}$ and so,

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\left\|\theta^{-1}\right\|_{1} \leqslant 2\left(n-\|\varphi\|_{1}\right)\left(\|\varphi\|_{1}-1\right) .
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Finally, the parabola $f(x)=2(n-x)(x-1)$ takes its maximum at $x=\frac{n+1}{2}$ and so,


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## The rank 2 case: $\alpha_{2}$

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For $n \geqslant n_{0}$ we have $\alpha_{2}(n) \geqslant \frac{n^{2}}{16}$.
So, the global known picture is

(v) $\beta_{r}(n) \leqslant K n^{M}$
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for some constants $K=K(r), M=M(r)$, and for $n \geqslant n_{0}$.

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(i) $\frac{n^{2}}{16} \leqslant \alpha_{2}(n) \leqslant \frac{(n-1)^{2}}{2}$,
(ii) $\beta_{2}(n)=n$,
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(iv) $K n^{r} \leqslant \alpha_{r}(n)$,
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## THANKS

