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# **Outline**

1. Introduction

- Introduction
- Strategy of the proof
- Orbit decidability
- Automaton groups

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Introduction

### Theorem (Sunic-V.)

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- Strategy of the proof

Will use results from Bogopolski-Martino-Ventura:

### Observation (B-M-V, 08)

Let H be f.g., and  $\Gamma \leqslant \operatorname{Aut}(H)$  f.g. If  $\Gamma \leqslant \operatorname{Aut}(H)$  is orbit undecidable then  $H \rtimes \Gamma$  has unsolvable CP.

and

## Proposition (B-M-V, 08)

For  $d \ge 4$ , there exist f.g., orbit undecidable, subgroups  $\Gamma \le GL_d(\mathbb{Z})$ .

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For  $d \ge 6$ , there exists a f.p. group G simultaneously satisfying the following three conditions:

- G is  $\mathbb{Z}^d$ -by-free,
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### Definition

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First examples:  $H = \mathbb{Z}^{q}$ 

### Observation (folklore

The full group  $\operatorname{Aut}(\mathbb{Z}^d) = \operatorname{GL}_d(\mathbb{Z})$  is orbit decidable.

**Proof.** For  $u, v \in \mathbb{Z}^d$ , there exists  $A \in GL_d(\mathbb{Z})$  such that v = Au if and only if  $gcd(u_1, \dots, u_d) = gcd(v_1, \dots, v_d)$ .

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Examples over the free group:  $H = F_r$ 

### Theorem (Whitehead'30)

The full group  $\operatorname{Aut}(F_r)$  is orbit decidable. That is, given  $u, v \in F_r$  one can decide whether  $v = \alpha(u)$  for some  $\alpha \in \operatorname{Aut}(F_r)$ .

Proof. This is a classical and very influential result.

### Theorem (Brinkmann, 06

Cyclic groups of  $\operatorname{Aut}(F_r)$  are orbit decidable. That is, given  $\varphi \in \operatorname{Aut}(F_r)$  and  $u, v \in F_r$ , one can decide whether  $v = \varphi^n(u)$ , up to conjugacy, for some  $n \in \mathbb{Z}$ .



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# Connection to semidirect products

### Observation (B-M-V)

Let H be f.g., and  $\Gamma \leqslant \operatorname{Aut}(H)$  f.g. If  $H \rtimes \Gamma$  has solvable CP, then  $\Gamma \leqslant \operatorname{Aut}(H)$  is orbit decidable.

**Proof.**  $G = H \times \Gamma$  contains elements  $(h, \gamma) \in H \times \Gamma$  operated like

$$(h_1, \gamma_1) \cdot (h_2, \gamma_2) = (h_1 \gamma_1(h_2), \gamma_1 \gamma_2)$$

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For  $h_1, h_2 \in H \leqslant G$ , we have  $h_1 \sim_G h_2 \Leftrightarrow \exists (h, \gamma) \in H \rtimes \Gamma$  s.t.

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Hence,  $h_1 \sim_G h_2 \Leftrightarrow \exists \gamma \in \Gamma$  and  $h \in H$  s.t.  $h_1 = h_{\gamma}(h_2)h^{-1}$ .  $\square$ 



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In fact, for the free and free abelian cases (among others), the convers is also true, after "erasing the relations from  $\Gamma$ ":

3. Orbit decidability

## Theorem (B-M-V, 08)

Let H be  $\mathbb{Z}^d$  or  $F_r$ , and  $\Gamma \leq \operatorname{Aut}(H)$  generated by  $\alpha_1, \ldots, \alpha_m$ . Then,  $H \rtimes_{\alpha_1,...,\alpha_m} F_m$  has solvable CP if and only if  $\Gamma = \langle \alpha_1, \dots, \alpha_m \rangle \leqslant \operatorname{Aut}(H)$  is orbit decidable.

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If  $\Gamma = \langle M_1, \dots, M_m \rangle$  is of finite index in  $GL_d(\mathbb{Z})$  then  $\mathbb{Z}^d \rtimes_{M_1, \dots, M_m} F_m$  has solvable conjugacy problem.

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# Corollary (Bogopolski-Martino-Maslakova-V., 06)

Free-by-cyclic groups have solvable conjugacy problem.

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# Observation (B-M-V, 08)

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### Corollary (Bogopolski-Martino-Maslakova-V., 06)

Free-by-cyclic groups have solvable conjugacy problem.

### Corollary

If  $\Gamma = \langle \varphi_1, \dots, \varphi_m \rangle$  has finite index in  $\operatorname{Aut}(F_r)$  then  $F_r \rtimes_{\varphi_1, \dots, \varphi_m} F_m$  has solvable conjugacy problem.

### Corollary

Every F<sub>2</sub>-by-free group has solvable conjugacy problem.

What we shall use is:

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#### But...

#### Theorem (Miller, 70's

There are free-by-free groups with unsolvable conjugacy problem.

So, there must be orbit undecidable subgroups in Aut  $(F_r)$ , for  $r \ge 3$ . Where are them ?

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Let H be a group, and let  $A \leq B \leq \operatorname{Aut}(H)$  and  $v \in H$  be such that  $B \cap \operatorname{Stab}^*(v) = 1$ . Then,

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$$\{\phi \in B \mid v\phi = w\} = B \cap (Stab(v) \cdot \varphi) = (B \cap Stab(v)) \cdot \varphi = \{\varphi\}.$$

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So,...

Taking the copy B of  $F_2 \times F_2$  in Aut( $F_3$ ) via the embedding

$$\begin{array}{ccccccc} F_2 \times F_2 & \hookrightarrow & & \textit{Aut}(F_3), \\ (u,v) & \mapsto & _u\theta_v \colon F_3 & \to & F_3 \\ & q & \mapsto & u^{-1}qv \\ & a & \mapsto & a \\ & b & \mapsto & b \end{array}$$

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### Proposition (B-M-V, 08)

For  $d \ge 4$ , there exist f.g., orbit undecidable, subgroups  $\Gamma \le GL_d(\mathbb{Z})$ .

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**Proof.** Consider 
$$F_2\simeq \langle P=\left(\begin{array}{cc}1&1\\1&2\end{array}\right),\ Q=\left(\begin{array}{cc}2&1\\1&1\end{array}\right)
angle \leq_{24} GL_2(\mathbb{Z}).$$

- $Stab(1,0) = \{M \mid (1,0)M = (1,0)\} = \{\begin{pmatrix} 1 & 0 \\ n & +1 \end{pmatrix} \mid n \in \mathbb{Z}\}.$
- $\langle P, Q \rangle \cap Stab(1,0) = \langle \begin{pmatrix} 1 & 0 \\ 12 & 1 \end{pmatrix} \rangle$ .
- Choose a free subgroup  $F_2 \simeq \langle P', Q' \rangle < \langle P, Q \rangle$  such that  $\langle P', Q' \rangle \cap Stab(1,0) = \{I\}$  and consider

$$B = \langle \left( \begin{array}{c|c} P' & 0 \\ \hline 0 & I \end{array} \right), \, \left( \begin{array}{c|c} Q' & 0 \\ \hline 0 & I \end{array} \right), \, \left( \begin{array}{c|c} I & 0 \\ \hline 0 & P' \end{array} \right), \, \left( \begin{array}{c|c} I & 0 \\ \hline 0 & Q' \end{array} \right) \rangle \leq GL_4(\mathbb{Z}).$$

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### Corollary

For  $d \geqslant 4$ , there exist  $\mathbb{Z}^d$ -by-free groups with unsolvable conjugacy problem.

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- Write v = (1, 0, 1, 0). By construction,  $B \cap Stab(v) = \{I\}$ .
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# Playing with 2 extra dimensions...

These orbit undecidable examples  $\Gamma \leqslant GL_4(\mathbb{Z})$  come from Mihailova's construction, so they are not finitely presented...

#### Proposition (Sunic-V.

For  $d \ge 6$ ,  $GL_d(\mathbb{Z})$  contains f.g., orbit undecidable, free, subgroups

### **Proof.** Let $d \ge 6$

- Since  $d-2\geqslant 4$ , there exists  $\langle g_1,\ldots,g_m\rangle=\Gamma\leqslant \operatorname{GL}_{d-2}(\mathbb{Z})$  being orbit undecidable.
- Let  $F_m = \langle f_1, \dots, f_m \rangle$ , and choose matrices  $s_1, \dots, s_m \in GL_2(\mathbb{Z})$  such that  $\langle s_1, \dots, s_m \rangle \simeq F_m$ .
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- Easy to see that  $F \leqslant \operatorname{GL}_d(\mathbb{Z})$  is orbit undecidable (using the orbit undecidability of  $\langle g_1, \dots, g_m \rangle = \Gamma \leqslant \operatorname{GL}_{d-2}(\mathbb{Z})$ ).  $\square$

In summary,

For  $d \geqslant 6$ , there exists a free  $\Gamma \leqslant \operatorname{GL}_d(\mathbb{Z})$  such that  $\mathbb{Z}^d \rtimes \Gamma$  has unsolvable CP.

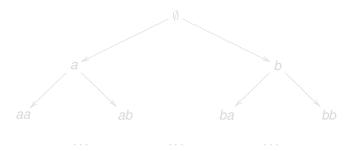
### Outline

- 1 Introduction
- Strategy of the proof
- Orbit decidability
- 4 Automaton groups

### Tree automorphisms

#### (joint work with Z. Sunic)

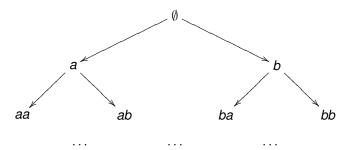
Let X be an alphabet on k letters, and let  $X^*$  be the free monoid on X, thought as a rooted k-ary tree:



### Tree automorphisms

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#### Definition

- A set of tree automorphisms is self-similar if it contains all sections of all of its elements.
- A finite automaton is a finite self-similar set (elements are called states).
- The group G(A) of tree automorphisms generated by an automaton A is called an automaton group.

#### Theorem (Sunic-V.)

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### Conclusion

So, we have proved that

#### Theorem

For  $d \geqslant 6$ , there exists  $\Gamma \leqslant GL_d(\mathbb{Z})$  free and orbit undecidable. Hence, the group  $\mathbb{Z}^d \rtimes \Gamma$ 

- is an automaton group,
- is free abelian-by-free,
- has unsolvable conjugacy problem.

# **THANKS**