# A recursive presentation for Mihailova's subgroup

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(joint with O. Bogopolski)



### Outline

- Mihailova's subgroup
- 2 Asphericity
- 3 The recursive presentation
- Orbit decidability

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Consider a finite presentation  $H = \langle x_1, \dots, x_n \mid R_1, \dots, R_m \rangle$ .

Associated to it K.A. Mihailova in 1958 considered

$$M(H) = \{(w_1, w_2) \in F_n \times F_n \mid w_1 =_H w_2\} \leqslant F_n \times F_n$$

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#### **Observation**

M(H) is generated by  $\{(x_1, x_1), \dots, (x_n, x_n), (1, R_1), \dots, (1, R_m)\}.$ 

- $W_1 =_H W_2$ 
  - $w_1^{-1}w_2$  is a product of conjugates of the  $R_j$ 's
  - For every  $z \in F_n$ ,  $(1, z^{-1}R_jz) = (z^{-1}, z^{-1})(1, R_j)(z, z)$
  - $(1, w_1^{-1} w_2) \in \langle (x_1, x_1), \dots, (x_n, x_n), (1, R_1), \dots, (1, R_m) \rangle$
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# Membership problem

#### Observation

 $MP(M(H), F_n \times F_n)$  is solvable  $\iff WP(H)$  is solvable

Proof. Obvious.

 $MP(M(H), F_n \times F_n)$ : given  $(w_1, w_2) \in F_n \times F_n$  decide whether  $(w_1, w_2) \in M(H)$  or not.

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- $F_n \times F_n$  has solvable word problem
- so, M(H) has solvable word problem
- so, M(H) is recursively presented.
- F. Grunewald, 1978: M(H) is finitely presented  $\iff$  H is finite.
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### Main result

#### Theorem (O. Boopolski, E.V.)

Let  $F_n$  be the free group on  $x_1, \ldots, x_n$ , and let  $H = \langle x_1, \ldots, x_n | R_1, \ldots, R_m \rangle$  be a finite, concise and Peiffer aspherical presentation. Then Mihailova's group  $M(H) \leqslant F_n \times F_n$  admits the following presentation

$$M(H) \simeq \langle d_1, \ldots, d_n, t_1, \ldots, t_m \mid [t_j, d^{-1}t_i^{-1}r_i d], [t_i, \operatorname{root}(r_i)] \rangle$$

 $(1 \leqslant i,j \leqslant m,\ d \in D_n)$ , where  $D_n$  is the free group with basis  $d_1,\ldots,d_n$ , and  $r_i$  denotes the word in  $D_n$  obtained from  $R_i$  by replacing each  $x_k$  to  $d_k$ .

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### Conciseness

Let  $H = \langle x_1, \dots, x_n | R_1, \dots, R_m \rangle$  be an arbitrary finite presentation (here,  $R_j$  are words on  $x_1, \dots, x_n$  which we'll assume reduced).

#### Definition

 $H = \langle x_1, \dots, x_n | R_1, \dots, R_m \rangle$  is concise if every  $R_j$  is non-trivial, and every two relations  $R_i$ ,  $R_j$ ,  $i \neq j$ , are not conjugate to each other or to the inverse of each other.

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# Identities among relations

#### Definition

Let  $U_1, \ldots, U_l \in F_n$ , and  $R_{i_1}, \ldots, R_{i_l} \in \{R_1, \ldots, R_n\}$ , and suppose

$$(U_1 R_{i_1}^{\varepsilon_1} U_1^{-1}) \cdots (U_l R_{i_l}^{\varepsilon_l} U_l^{-1}) = 1$$

holds in  $F_n$ . In this situation, the sequence of elements of  $F_n$ 

$$(U_1R_{i_1}^{\varepsilon_1}U_1^{-1},\ldots,U_lR_{i_l}^{\varepsilon_l}U_l^{-1})$$

is called an identity among relations of length I.

For I = 0 we have the empty identity among relations, ( ).

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### Peiffer transformations

Define the following elementary transformations of an identity among relations  $(U_1 R_i^{\varepsilon_1} U_1^{-1}, \dots, U_l R_i^{\varepsilon_l} U_l^{-1})$ :

- deletion/insertion: if  $(U_p R_{i_p}^{\varepsilon_p} U_p^{-1}) \cdot (U_{p+1} R_{i_{p+1}}^{\varepsilon_{p+1}} U_{p+1}^{-1}) = 1$ , delete them.
- exchange: replace two consecutive terms, say

$$U_p R_{i_p}^{\varepsilon_p} U_p^{-1}$$
 and  $U_{p+1} R_{i_{p+1}}^{\varepsilon_{p+1}} U_{p+1}^{-1}$ ,

by the new ones

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# Peiffer asphericity

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The presentation  $H = \langle x_1, \dots, x_n | R_1, \dots, R_m \rangle$  is Peiffer aspherical if every identity among relations can be carried to the empty one by a sequence of Peiffer transformations.

 Chiswell, D. Collins, J. Huebschmann (1981): asphericity is preserved under free products, certain HNN extensions, and Tietze transformations.

In the literature other related concepts (diagrammatical asphericity, topological asphericity, ...).

#### Theorem (Collins, Miller, 1999)

There exists a finite, concise and Peiffer aspherical presentation  $H = \langle x_1, \dots, x_n | R_1, \dots, R_m \rangle$  with unsolvable word problem.



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$$\begin{array}{ccc}
\pi \colon F_{n+m} & \to & M(H) \\
d_i & \mapsto & (x_i, x_i) \\
t_j & \mapsto & (1, R_j).
\end{array}$$

- $\pi$  is clearly onto,
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to every  $w \in \ker(\pi)$  we'll associate an identity among relations iar(w) such that

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Let  $w \in \ker(\pi)$  and write it as  $w = u_1 t_{i_1}^{\varepsilon_1} u_2 \cdots u_l t_{i_l}^{\varepsilon_l} u_{l+1}$  (where  $l \geqslant 0$  and  $u_1, \dots, u_{l+1}$  are words in the  $d_i's$ ).

Projecting  $\pi(w)$  to each coordinate, we have

$$U_1 U_2 \cdots U_{l+1} = 1$$
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Now, denote the accumulative products by  $\mathbb{U}_i = U_1 U_2 \cdots U_i$ ,  $i = 1, \dots, l+1$  (note that  $\mathbb{U}_{l+1} = 1$ ), and we have

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## Outline

- Mihailova's subgroup
- Asphericity
- The recursive presentation
- Orbit decidability

Let

$$1 \longrightarrow F \stackrel{\alpha}{\longrightarrow} G \stackrel{\beta}{\longrightarrow} H \longrightarrow 1$$

### be an algorithmic short exact sequence of groups such that

- (i) TCP(F) is solvable
- (ii) CP(H) is solvable
- (iii) there is an algorithm which, given an input  $1 \neq h \in H$ , computes a finite set of elements  $z_{h,1}, \ldots, z_{h,t_h} \in H$  such that

$$C_H(h) = \langle h \rangle z_{h,1} \sqcup \cdots \sqcup \langle h \rangle z_{h,t_h}.$$

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Particularizing to the case where *F* and *H* are free, we obtain:

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#### Theorem (Miller, 70's

There are free-by-free groups with unsolvable conjugacy problem.

So, there must be orbit undecidable subgroups in Aut  $(F_n)$ , for  $n \ge 3$ . Where are them ?

### Proposition (Bogopolski-Martino-V., 2008)

Let F be a group, and let  $A \leq B \leq Aut(F)$  and  $w \in F$  be such that  $B \cap Stab^*(w) = 1$ . Then,

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### Corollary

Let F be a group, and let  $A \leq B \leq \operatorname{Aut}(F)$  and  $w \in F$  be such that  $B \cap \operatorname{Stab}^*(w) = 1$ . If  $B \simeq F_2 \times F_2$  and A is the Mihailova subgroup corresponding to a group with unsolvable word problem then,  $A \leq \operatorname{Aut}(F)$  is orbit undecidable.

Well, with the embedding 
$$F_2 \times F_2 \longrightarrow \operatorname{Aut}(F_3)$$
  $(u,v) \mapsto u\theta_v \colon F_3 \longrightarrow F_3$   $q \mapsto u^{-1}qv$   $a \mapsto a$   $b \mapsto b$ 

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Take  $F = \mathbb{Z}^n$  with  $n \geqslant 4$ .

 $F_2 \times F_2 \leqslant GL_2(\mathbb{Z}) \times GL_2(\mathbb{Z}) \leqslant GL_4(\mathbb{Z}) \leqslant GL_n(\mathbb{Z})$  (and  $TCP(\mathbb{Z}^4)$  is solvable).

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## **THANKS**