On the difficulty of inverting automorphisms of free groups

Enric Ventura

Departament de Matemàtica Aplicada III Universitat Politècnica de Catalunya

Groups in Galway, 2011

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Outline

- Motivation
- 2 Free groups
- 3 Lower bounds: a good enough example
- 4 Upper bounds: outer space
- 5 The special case of rank 2

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(Joint work with P. Silva and M. Ladra.)

Find a group G where \cdot is "easy" but ()⁻¹ is "difficult".

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\phi \psi : F_3 \rightarrow F_3
         a \mapsto bc^{-1}a^{-1}bc
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$$F_{3} = \langle \textbf{\textit{a}}, \textbf{\textit{b}}, \textbf{\textit{c}} \mid \rangle.$$

$$\phi \colon F_{3} \to F_{3} \qquad \psi \colon F_{3} \to F_{3}$$

$$a \mapsto ab \qquad a \mapsto bc^{-1}$$

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$$F_{5} = \langle a, b, c, d, e \mid \rangle.$$

$$\psi_{n} \colon F_{5} \to F_{5} \qquad \psi_{n}^{-1} \colon F_{4} \to F_{4}$$

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$$e \mapsto d^{n}e \qquad e \mapsto (d^{-1}((b^{-1}a^{n})^{n}c)^{n})^{n}e.$$

- We have formalized the situation.
- We have seen that inverting in $Aut(F_r)$ is not that bad.
- We now want to look for worse groups G.



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Definition

Let $A = \{a_1, \dots, a_r\}$ be a finite alphabet, and $G = \langle A \mid R \rangle$ be a finite presentation for a group G. We have the word metric:

for
$$g \in G$$
, $|g| = \min\{n \mid g = a_{i_1}^{\epsilon_1} \cdots a_{i_n}^{\epsilon_n}\}$.

Definition

For $\theta \in Aut(G)$, note θ is determined by $a_1\theta, \ldots, a_r\theta$ and define

$$||\theta||_1 = |a_1\theta| + \cdots + |a_r\theta|,$$

$$||\theta||_{\infty} = \max\{|a_1\theta|,\ldots,|a_r\theta|\}.$$

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Let $G = \langle A \mid R \rangle$ be a finite presentation for G. We define the function:

$$\alpha_{A}(n) = \max\{||\theta^{-1}||_{1} \mid \theta \in Aut(G), ||\theta||_{1} \leqslant n\}.$$

Clearly,
$$\alpha_A(n) \leqslant \alpha_A(n+1)$$
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The bigger is α_A , the more "difficult" will be to invert automorphisms of G (with respect to the given set of generators A).

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Determine the asymptotic growth of the function α_A .

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For the rest of the talk, $G = F_r = \langle a_1, \dots, a_r \mid \rangle$.

Definition

Every $w \in F_r$ has its length, |w|, and its cyclic length, |w|: $|a_1a_1^{-1}a_2| = |a_2| = |a_2| = 1$, $|a_1a_2a_1^{-2}| = 4$, $|a_1a_2a_1^{-2}| = |a_2a_1^{-1}| = 2$.

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i) $|w^n| \le |n||w|$ and $|w^n| = |n||w|$; ii) $|vw| \le |v| + |w|$, but $|vw| \le |v| + |w|$ is not true in general.

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Are these functions equal up to multiplicative constants?

 α_r and γ_r are not; β_r is not clear.

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For rank r = 2 we have

(i) for
$$n \ge 4$$
, $\alpha_2(n) \le \frac{(n-1)^2}{2}$,

(ii) for
$$n \geqslant n_0$$
, $\alpha_2(n) \geqslant \frac{n^2}{16}$

(iii) for
$$n \geqslant 1$$
, $\beta_2(n) = n$,

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Theorem

For $r \geqslant 3$ there exist K = K(r) and M = M(r) such that, for $n \geqslant 1$,

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Proof: For $r \ge 2$ and $n \ge 1$, consider

A straightforward calculation shows that

$$\|\psi_{r,n}\|_1 = \|\psi_{r,n}\|_1 = (r-1)n + r$$
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Hence, for $n \ge r$,

$$\gamma_r(rn) \geqslant \gamma_r((r-1)n+r) \geqslant n^{r-1}.$$

Now, for *n* big enough, take the closest multiple of *r* below,

$$n \geqslant rm > n - r$$
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and

$$\gamma_r(n)\geqslant \gamma_r(rm)\geqslant m^{r-1}>\left(\frac{n-r}{r}\right)^{r-1}=\left(\frac{n}{r}-1\right)^{r-1}\geqslant \frac{1}{2r^{r-1}}n^{r-1}.\quad \Box$$

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A lower bound for α_r

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Outline

- Motivation
- 2 Free groups
- 3 Lower bounds: a good enough example
- 4 Upper bounds: outer space
- 5 The special case of rank 2

To prove the upper bound

(ii)
$$\beta_r(n) \leqslant Kn^M$$
,

we'll need to use the recently discovered metric in the outer space \mathcal{X}_r .

- By graf Γ we mean a finite, connected graph of rank r, with no vertices of degree 1 or 2.
- A metric on Γ is a map $\ell \colon E\Gamma \to [0,1]$ such that $\sum_{e \in E\Gamma} \ell(e) = 1$, and $\{e \in E\Gamma \mid \ell(e) = 0\}$ is a forest.
- For a graph Γ, Σ_Γ = {metrics on Γ} = a simplex with missing faces.
- If $\Gamma' = \Gamma/$ forest, then we identify points in $\Sigma_{\Gamma'}$ with the corresponding points in Σ_{Γ} by assigning length 0 to the collapsed edges.
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$$d(x, y) \geqslant 0$$
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$$d(x,z) \leqslant d(x,y) + d(y,z)$$
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For any $\epsilon > 0$ there is constant $M = M(r, \epsilon)$ such that for all $x, y \in \mathcal{X}_r(\epsilon)$,

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$$\begin{array}{ll} \textit{d}(\textit{x},\,\phi\cdot\textit{x}) &=& \min\{\log(\sigma(\alpha)) \mid \alpha \text{ diff. markings}\}\\ &=& \log\big(\min\{\sigma(\phi\gamma_w\gamma_p) \mid \textit{w} \in \textit{F}_r,\, p = \text{ "half petal"}\}\big)\\ &\sim& \log\big(\min\{\sigma(\phi\gamma_w) \mid \textit{w} \in \textit{F}_r\}\big)\\ &=& \log\big(\min\{||\phi\gamma_w||_\infty \mid \textit{w} \in \textit{F}_r\}\big)\\ &=& \log(|||\phi|||_\infty)\\ &\sim& \log(|||\phi|||_1). \end{array}$$

Now, using Bestvina-AlgomKfir theorem,

$$\log(|||\phi^{-1}|||_1) = d(x, \phi^{-1} \cdot x) = d(\phi \cdot x, x) \leqslant Md(x, \phi \cdot x) = M \log(|||\phi|||_1).$$



Outline

- Motivation
- 2 Free groups
- 3 Lower bounds: a good enough example
- Upper bounds: outer space
- The special case of rank 2

The rank 2 case

These functions for $Aut(F_2)$ are much easier to understand due to the following technical lemmas.

Lemma

Let $\varphi \in Aut(F_2)$ be positive. Then φ^{-1} is cyclically reduced and $||\varphi^{-1}||_1 = ||\varphi||_1$.

Lemma

For every $\theta \in Aut(F_2)$, there exist two letter permuting autos $\psi_1, \ \psi_2 \in Aut(F_2)$, a positive one $\varphi \in Aut^+(F_2)$, and an element $g \in F_2$, such that $\theta = \psi_1 \varphi \psi_2 \lambda_{\sigma}$ and $||\varphi||_1 + 2|g| \leqslant ||\theta||_1$.

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Theorem

For every
$$\theta \in Aut(F_2)$$
, $||\theta^{-1}||_1 = ||\theta||_1$. Hence, $\gamma_2(n) = n$.

Proof. Let $\theta \in \text{Aut}(F_2)$, decomposed as above, $\theta = \psi_1 \varphi \psi_2 \lambda_g$. Then

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For
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 we have $\alpha_2(n) \leqslant \frac{(n-1)^2}{2}$.

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For $n \geqslant n_0$ we have $\alpha_2(n) \geqslant \frac{n^2}{16}$.

So, the global known picture is

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,

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(v)
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THANKS